

# Page Rank

(Section 3.3 - kind of)

## Google's Original Search Algorithm

In a paper by Brin & Page

(the founders of Google) from

1998. Since been improved

from its original framework.

Step 1: Make an  $n \times n$  matrix  $A$  where  
 $n =$  the number of webpages (urls)  
in existence ( $\approx 1.5$  billion),

$$A = (A_{i,k})_{i,k=1}^n \text{ where}$$

$$A_{i,k} = \begin{cases} 1, & \text{if page } k \text{ links to} \\ & \text{page } i \\ 0, & \text{otherwise} \end{cases}$$

By convention, a page doesn't link to  
itself.

Example 1: (3 pages) Given pages

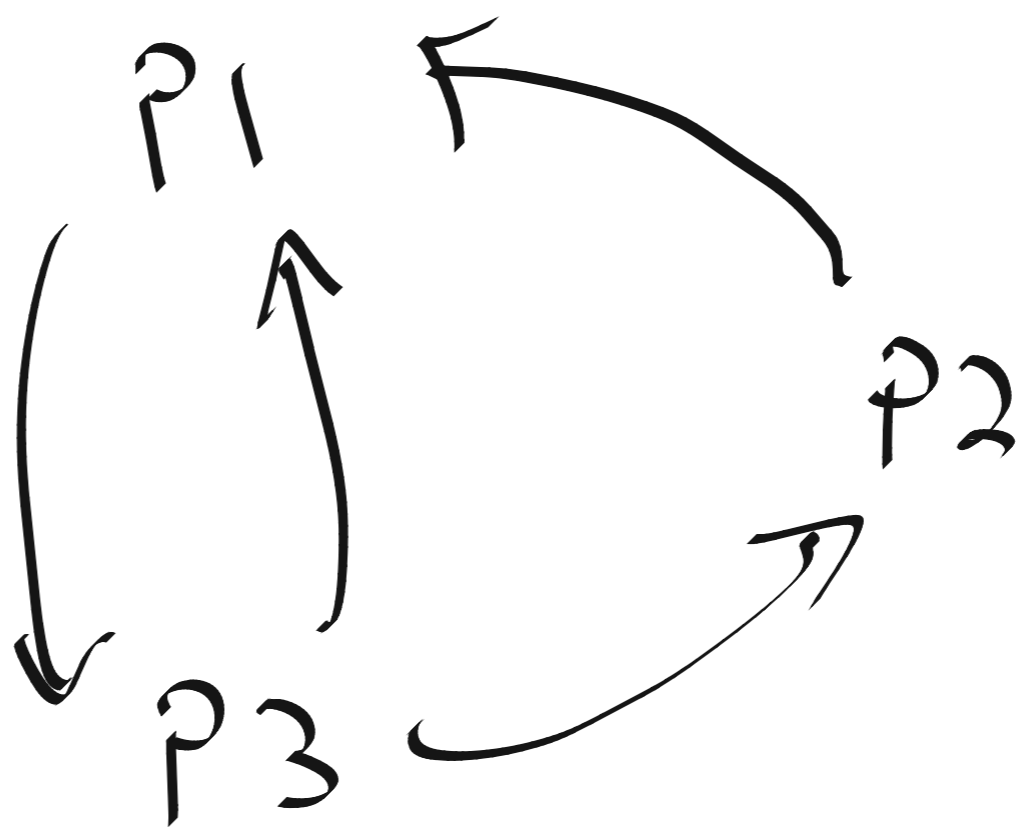
P1, P2, and P3, suppose

P1 links to P3

P2 links to P1

P3 links to P1 and P2

Picture:



Then

$$A^{-1} = \begin{bmatrix} & p_1 & & & & \\ & & p_2 & & & \\ & & & p_3 & & \\ 0 & & & & & \\ 0 & & & & & \\ 1 & & & & & \end{bmatrix}$$

Step 2: Add the entries of the  $k^{\text{th}}$  column of  $A$  for all  $1 \leq k \leq n$  and divide the  $A_{i,k}$  entries of  $A$  by this number. Call this new matrix  $B$ .

Back to Example 1:

$$A = \begin{bmatrix} 0 & 1 & 1 \\ 0 & 0 & 1 \\ 1 & 0 & 0 \end{bmatrix}$$

Sum of entries in column 1 = 1

Sum of entries in column 2 = 1

Sum of entries in column 3 = 2

Can leave entries in columns 1 and 2 alone, divide the entries in column 3 by 2 to get  $B$  :

$$B = \begin{bmatrix} 0 & 1 & \frac{1}{2} \\ 0 & 0 & \frac{1}{2} \\ 1 & 0 & 0 \end{bmatrix}$$

Note: the entries in every column of  $B$  add up to one.

**Problem:** What if all the entries in one column are zero? Then we're dividing by zero...

**Solution:** Change all the entries in that column to ones. When constructing  $B$ , the entries will now be  $1/n$  where  $n =$  number of webpages.

As  $n \rightarrow \infty$ ,  $1/n \rightarrow 0$ ,

so this is (somewhat) representative of the original matrix if  $n$  is large.

**Step 3:** Build in the possibility that people will randomly start their search over with probability  $1-d$  (usually  $d = .85$ ) and let  $C$  be the  $n \times n$  matrix with entries

$$C_{ik} = d B_{i,k} + \frac{(1-d)}{n}$$

for all  $1 \leq i, k \leq n$ .

**Note:** every column of  $C$  still sums to one.



## Back to Example 1

Use fractions:  $d = .85 = \frac{85}{100} = \frac{17}{20}$

$$1-d = .15 = \frac{3}{20}$$

**Why?** Sometimes Wolfram Alpha will make mistakes with decimals.

$$C_{i,k} = \frac{17}{20} B_{i,k} + \frac{\frac{3}{20}}{3}$$

$$C_{i,k} = \frac{17}{20} B_{i,k} + \frac{1}{20}$$

$$1 \leq i, k \leq 3$$

Can write this using matrices:

$$C = \frac{17}{20} B + \frac{1}{20} \begin{bmatrix} 1 & 1 & 1 \\ 1 & 1 & 1 \\ 1 & 1 & 1 \end{bmatrix}$$

Recall  $B = \begin{bmatrix} 0 & 1 & 1/2 \\ 0 & 0 & 1/2 \\ 1 & 0 & 0 \end{bmatrix}$ , so

$$C = \frac{1}{20} \left( 17B + \begin{bmatrix} 1 & 1 & 1 \\ 1 & 1 & 1 \\ 1 & 1 & 1 \end{bmatrix} \right)$$

$$= \frac{1}{20} \left( \begin{bmatrix} 0 & 17 & 17/2 \\ 0 & 0 & 17/2 \\ 17 & 0 & 0 \end{bmatrix} + \begin{bmatrix} 1 & 1 & 1 \\ 1 & 1 & 1 \\ 1 & 1 & 1 \end{bmatrix} \right)$$

$$= \frac{1}{20} \begin{bmatrix} 1 & 18 & 19/2 \\ 1 & 1 & 19/2 \\ 18 & 1 & 1 \end{bmatrix}$$

(can stop here)

$$= \begin{bmatrix} 1/20 & 18/20 & 19/40 \\ 1/20 & 1/20 & 19/40 \\ 18/20 & 1/20 & 1/20 \end{bmatrix}$$

Observe that every column  
still adds up to one!

Step 4:  $C$  will always have  $\lambda = 1$  as its largest eigenvalue.

Take the corresponding eigenvector  $\vec{v}$  and divide every entry of  $\vec{v}$  by

$$\|\vec{v}\|_1 = \sum_{i=1}^n |v_i|.$$

The  $i$ th entry of this vector is the PageRank of page  $P_i$  for all  $1 \leq i \leq n$ .

Back to Example 1:

$$C = \frac{1}{20} \begin{bmatrix} 1 & 18 & 19/2 \\ 1 & 1 & 19/2 \\ 18 & 1 & 1 \end{bmatrix}$$

Eigenvalue  $\lambda = 1$  has  
associated eigenvector

$$\vec{v} = \begin{bmatrix} 703/686 \\ 190/343 \\ 1 \end{bmatrix} \quad (\text{entries are all positive})$$

$$\|\vec{v}\|_1 = \frac{703}{686} + \frac{190}{343} + 1$$

$$\|\vec{v}\|_1 = \frac{1769}{686}$$

Divide  $\vec{v}$  by this number.

$$\frac{\vec{v}}{\|\vec{v}\|} = \frac{686}{1769} \left[ \begin{array}{l} 703/686 \\ 190/343 \\ 1 \end{array} \right]$$

$$= \left[ \begin{array}{l} 703/1769 \\ 380/1769 \\ 686/1769 \end{array} \right]$$

Observe that the entries of  $\frac{\vec{v}}{\|\vec{v}\|}$ ,

sum to one:

$$\frac{703}{1769} + \frac{380}{1769} + \frac{686}{1769} = \frac{1769}{1769} = 1$$

Each entry represents the probability of landing at the associated page, called the PageRank.

$$\frac{\vec{v}}{\|\vec{v}\|_1} = \begin{bmatrix} 703/1769 \\ 380/1769 \\ 686/1769 \end{bmatrix} \begin{array}{l} \leftarrow \text{PageRank of } P1 \\ \leftarrow \text{PageRank of } P2 \\ \leftarrow \text{PageRank of } P3 \end{array}$$

Example 2: (four pages) find the PageRank  
of  $P_2$  if

$P_1$  links to  $P_2$  and  $P_4$

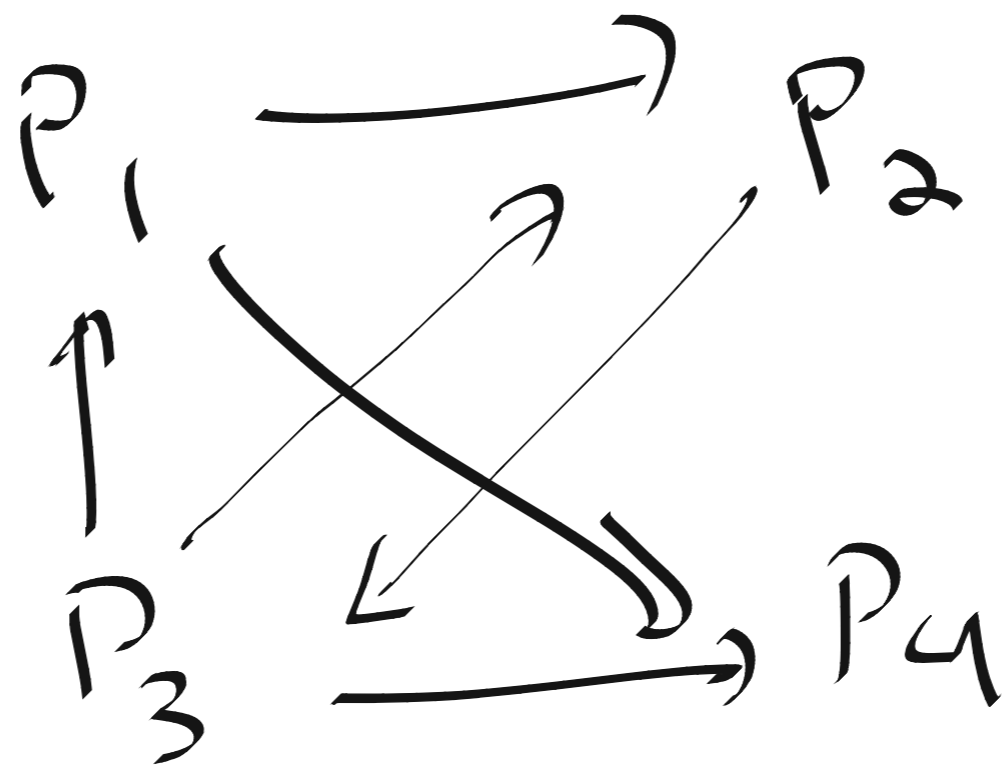
$P_2$  links to  $P_3$

$P_3$  links to  $P_1, P_2,$  and  $P_4$

$P_4$  links to nothing.

Solution:

Picture





Step 1: Make A.

$$A = \begin{array}{c} P_1 \quad P_2 \quad P_3 \quad P_4 \\ \left[ \begin{array}{cccc} 0 & 0 & 1 & 0 \\ 1 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 \\ 1 & 0 & 1 & 0 \end{array} \right] \end{array}$$

Step 2: Make B. Note column 4 is all zeros, so change the zeros to ones in this

column:

$$A \rightarrow \left[ \begin{array}{cccc} 0 & 0 & 1 & 1 \\ 1 & 0 & 1 & 1 \\ 0 & 1 & 0 & 1 \\ 1 & 0 & 1 & 1 \end{array} \right]$$

$$B = \begin{bmatrix} 0 & 0 & \frac{1}{3} & \frac{1}{4} \\ \frac{1}{2} & 0 & \frac{1}{3} & \frac{1}{4} \\ 0 & 1 & 0 & \frac{1}{4} \\ \frac{1}{2} & 0 & \frac{1}{3} & \frac{1}{4} \end{bmatrix}$$

Note every column of  $B$  adds up to one.

Step 3: Make  $C$ .

$$d = .85 = \frac{17}{20}$$

$$C_{i,k} = \frac{17}{20} B_{i,k} + \frac{\frac{3}{20}}{4}$$

$$C_{i,k} = \frac{17}{20} B_{i,k} + \frac{3}{80}$$

$$1 \leq i, k \leq 4.$$

$$C = \frac{17}{20} B + \frac{3}{80} \begin{bmatrix} | & | & | & | \\ | & | & | & | \\ | & | & | & | \\ | & | & | & | \end{bmatrix}$$

$$= \frac{1}{20} \left( 17B + \frac{3}{4} \begin{bmatrix} | & | & | & | \\ | & | & | & | \\ | & | & | & | \\ | & | & | & | \end{bmatrix} \right)$$

$$= \frac{1}{20} \left( \begin{bmatrix} 0 & 0 & 17/3 & 17/4 \\ 17/2 & 0 & 17/3 & 17/4 \\ 0 & 17 & 0 & 17/4 \\ 17/2 & 0 & 17/3 & 17/4 \end{bmatrix} + \right.$$

$$\left. \frac{3}{4} \begin{bmatrix} | & | & | & | \\ | & | & | & | \\ | & | & | & | \\ | & | & | & | \end{bmatrix} \right)$$

$$= \frac{1}{20} \begin{bmatrix} 3/4 & 3/4 & 77/12 & 5 \\ 37/4 & 3/4 & 77/12 & 5 \\ 3/4 & 7/4 & 3/4 & 5 \\ 37/4 & 3/4 & 77/12 & 5 \end{bmatrix}$$

Observe each column of  $C$   
adds up to one.

**Step 4:** Compute PageRanks by finding  
the eigenvector associated to  $\lambda = 1$   
for  $C$ .

We get the eigenvector

$$\vec{v} = \begin{bmatrix} 40/57 \\ 1 \\ 1769/1463 \\ 1 \end{bmatrix}$$

$$\|\vec{v}\|_1 = 40/57 + 1 + \frac{1769}{1463} + 1$$

$$\|\vec{v}\|_1 = \frac{17,165}{4389}$$

$$\frac{\vec{v}}{\|\vec{v}\|_1} = \frac{4389}{17,165} \left[ \begin{array}{c} 40/57 \\ 1 \\ 1769/1463 \\ 1 \end{array} \right]$$

Only asked for the PageRank of

$P_2$ , which is

$$\frac{4389}{17,165}$$

For large values of  $n$ , how

Page and Brin find this eigenvector

is Google's secret!

Some questions:

1) How do we know  $\lambda=1$  is always an eigenvalue for  $C$ ?

2) How do we know  $\lambda=1$  is the largest eigenvalue?

3) The associated eigenvector  $\vec{v}$  had better have positive entries.

How do we know this is the case?

Answer for 1) : Eigenvalues of  $C =$   
eigenvalues of  $C^t$ , and

$$\vec{w} = \begin{bmatrix} | \\ | \\ | \\ | \\ | \end{bmatrix} \text{ is}$$

an associated eigenvector  
for  $\lambda = 1$  and  $C^t$ .

If we look at (for any  $n \times n$  matrix  $A$ )

$$\det(A^t - \lambda I_n)$$

$$= \det(A^t - \lambda I_n^t)$$

$$= \det((A - \lambda I_n)^t)$$

$$= \det(A - \lambda I_n) \quad (\text{determinant property 4})$$

So  $0 = \det(A^t - \lambda I_n) = \det(A - \lambda I_n)$ ,

which tells us that the eigenvalues of  $A$  and  $A^t$  are the same.

Now consider  $C^t$ . Remember that

$$(C^t)_{i,k} = C_{k,i} \text{ for all } (1 \leq i, k \leq n)$$

$$\text{and that } \sum_{k=1}^n C_{k,i} = 1 \text{ for all}$$

fixed  $i$ ,  $1 \leq i \leq n$ .



Then

$$C^t \begin{bmatrix} \vdots \\ \vdots \\ \vdots \end{bmatrix}$$

$$= \begin{bmatrix} \sum_{k=1}^n (C^t)_{1,k} \\ \sum_{k=1}^n (C^t)_{2,k} \\ \vdots \\ \sum_{k=1}^n (C^t)_{n,k} \end{bmatrix}$$

$$= \begin{bmatrix} \sum_{k=1}^n C_{k,1} \\ \sum_{k=1}^n C_{k,2} \\ \vdots \\ \sum_{k=1}^n C_{k,n} \end{bmatrix} \begin{matrix} = 1 \\ = 1 \\ \vdots \\ = 1 \end{matrix} \quad \begin{matrix} \text{(all} \\ \text{columns} \\ \text{of } C \\ \text{sum to} \\ \text{1)} \end{matrix}$$

$$= \begin{bmatrix} | \\ | \\ | \\ | \\ | \end{bmatrix}$$

$$= 1 \cdot \begin{bmatrix} | \\ | \\ | \\ | \\ | \end{bmatrix}$$

This shows  $\lambda=1$  is an  
eigenvalue for  $C^t$ , so  $\lambda=1$   
is an eigenvalue for  $C$  but  
not with the vector of all ones!

Answer for 2): norm inequality, takes a  
little bit of work

Answer for 3): Perron-Frobenius  
Theorem!